

Chapter 10. Synchronous Serial Communication and Keyboard Connection

1. Synchronous Communication

As its name implies, synchronous communication takes place between a transmitter and a receiver operating on synchronized clocks. In a synchronous system, the communication partners have a short conversation before data exchange begins. In this conversation, they align their clocks and agree upon the parameters of the data transfer, including the time interval between bits of data. Any data that falls outside these parameters will be assumed to be either in error or a placeholder used to maintain synchronization. (Synchronous lines must remain constantly active in order to maintain synchronization, thus the need for placeholders between valid data.) Once each side knows what to expect of the other, and knows how to indicate to the other whether what was expected was received, then communication of any length can commence.

Even though 16F877's USART module provides hardware enabled synchronous master/slave mode of serial communication, we opt to a software enabled approach. It's because the built-in serial port will be connected to a host PC for hex code download. Of course, that same port can be used for other serial device, it would be cumbersome to connect and disconnect a code. Moreover, we will connect another serial device, like a keyboard or mouse, in the example of this chapter, therefore, software approach will give us more freedom of adding additional serial device.

An application of this chapter is to connect a keyboard (eventually two keyboards) and one LCD to the 16F877 in order to let two persons of hearing or speaking disability communicate by typing and reading. The keyboard we are going to connect is the most common type, IBM AT or PS/2 keyboard. The keyboard communicates with PC in synchronous serial communication.

AT type keyboard has 5 pins while PS/2 type keyboard has 6 pins. As illustrated below, for both types of keyboard, there are total 4 signals: +5V power signal, ground, and a CLOCK line, and a DATA line.

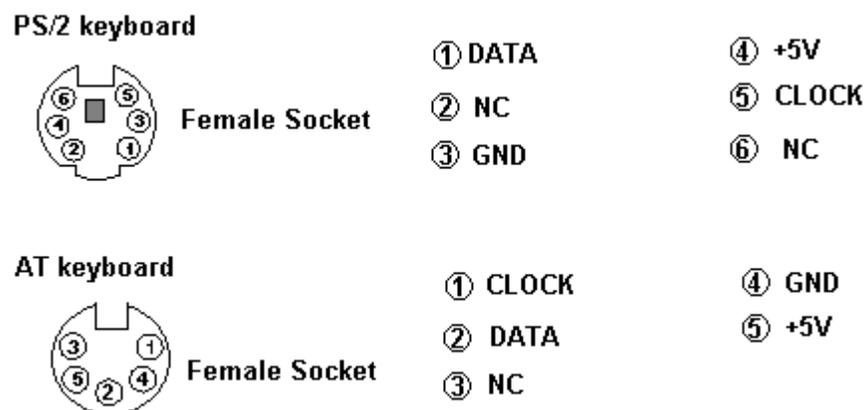


Fig. 75 PS/2 and AT type keyboard

2. IBM AT- or PS/2 – type Keyboard Protocol

The protocol between the keyboard and PC is the most important subject we have to understand in our example application.

The PS/2 mouse and keyboard implement a bidirectional synchronous serial protocol. The bus is "idle" when both lines are high (open-collector). This is the only state where the keyboard/mouse is allowed begin transmitting data. The host has ultimate control over the bus and may inhibit communication at any time by pulling the Clock line low.

The keyboard always generates the clock signal. This is done by a keyboard controlling microcontroller inside the keyboard.

The Data and Clock lines are both open-collector with pull-up resistors to +5V. An "open-collector" interface has two possible state: low, or high impedance. In the "low" state, a transistor pulls the line to ground level. In the "high impedance" state, the interface acts as an open circuit and doesn't drive the line low or high.

The first thing we have to know is how the keyboard controller chip send data to a host (PC or 17F877 in our case). As mentioned above there are two signals from the keyboard: CLOCK and DATA. DATA is sent only when synchronized with the CLOCK. When the keyboard is idle, without any key pressed, both CLOCK and DATA are remained pulled up High. When a key is pressed in the keyboard, both the CLOCK pulse and DATA pulse are transmitted from the keyboard. Through the DATA, strings of byte data are generated. The clock pulses are generated during the data transmission through the CLOCK line.

One thing we have to remember is that a single key stroke does not generate only a byte of data: it generates usually 3 bytes of data and, but other keys generate 5 bytes of data. The list of byte data generated by each individual key is called Keyboard Scan Codes. This discussion follows.

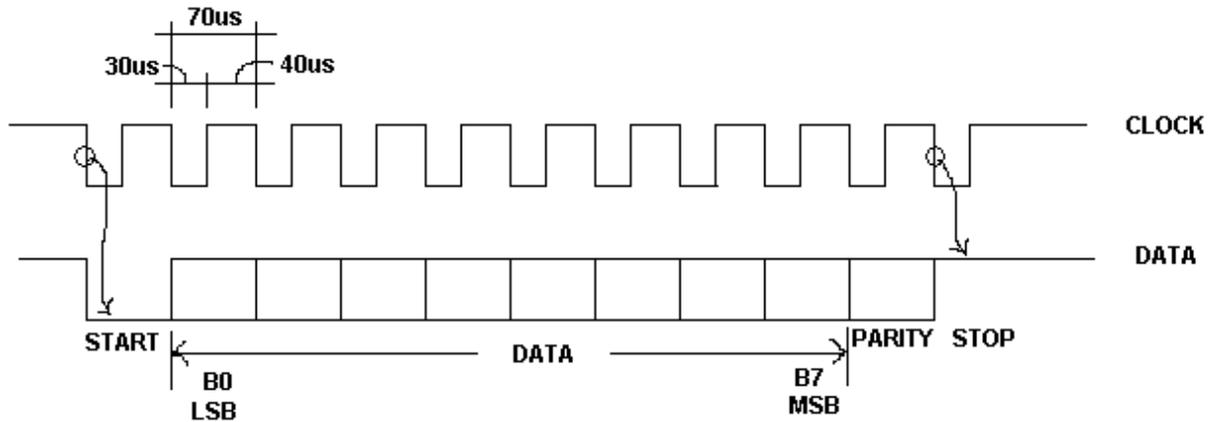
Let's continue our discussion on keyboard protocol. A byte data from the keyboard is sent in a frame consisting of 11 bits. The frame consists, in the following order, of:

- 1 Start bit (Low),
- 8-bit data (LSB first, as usual),
- 1 Odd Parity bit, and
- 1 Stop bit (High).

The width of the data bit is about 70 μ s. The frame is synchronized with 11 clock pulses of 70 μ s with about 40% duty cycle. Namely, the clock pulse's width is 70 μ s and it's High is about 307 μ s and its Low for 40 μ s. As indicated below, a host can sample (or monitor) each bit of the frame at the falling edge of the clock pulse. If we allow a short transition time of High-to-Low change, it would be safe to sample after around 5 μ s of the High-to-Low transition of the clock.

To read the frame using 16F877, we need two I/O ports configured as inputs for CLOCK and DATA lines. First we monitor the CLOCK line for transition from High to Low. When it

changes to Low, after 5 μ s, we read the DATA line for each bit. Once a bit is read, now we back to the CLOCK monitoring. The CLOCK must go back to High and do the High-to-Low transition for the next bit reading. This process goes for all 11 bits of a frame. Since the 8-bit byte is sent LSB first, as soon as each bit of the byte is read, it must be rotate to the right by one to make a regular byte format: MSB to LSB.



Before we further proceed to the Keyboard Scan Codes, let's have a 16F877 connection with a standard AT or PS/2 keyboard. CLOCK line is connected to RB7 and DATA line to RB6 as illustrated below.

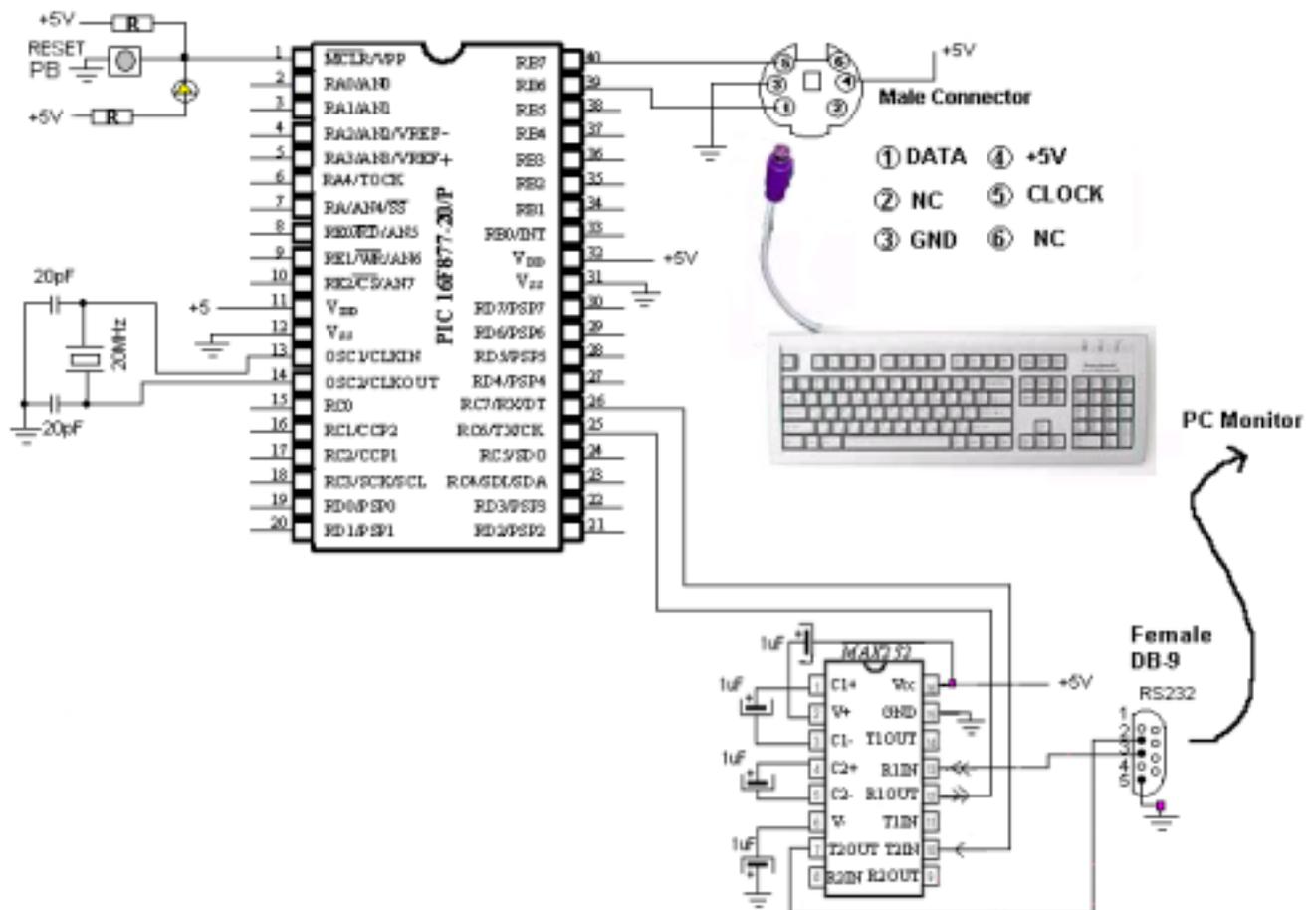


Fig. 76 PIC 16F877 connection to PS/2 Keyboard

Then, let's build a subroutine to read the 11-bit frame: the basic building block of keyboard reading. As explained above, reading each bit is based on the monitoring of the CLOCK line (RB7) of the High-to-Low transition. Since there are chances that the monitoring would be in the middle of the frame transmission, we may want to have a short High CLOCK before we allow to receive a frame. The subroutine, RX11bit, is listed below with ample amount of comments. In the subroutine, we keep the Parity bit for later use of transmission parity check. Also, an unknown transmission error is recorded when the last bit (STOP) is not High. The 8-bit data, after the subroutine, is stored in the file register of DATAreg.

```

;SUBROUTINE RX11bit =====
;RX Routine for 11-bit frame read
;1 Start
;8 Data (LSB first)
;1 Parity (Odd)
;1 Stop (HIGH)
;KSTAT Bit Info: KSTAT<0> : parity   KSTAT<2>:KBD Error
;
RX11bit
    clrf        DATAreg
    banksel    PORTB
;Let it have at least 200us CLOCK high period
    btfss     PORTB, CLOCK
    goto      RX11bit           ;if CLOCK is LOW, start again
    call     Delay100us        ;200uS delays
    call     Delay100us
;check again for CLOCK
    btfss     PORTB, CLOCK
    goto      RX11bit
;READY TO MONITOR CLOCK of H-to-L TRANSITION
Scheck
    btfsc     PORTB, CLOCK
    goto      Scheck
;CLOCK pulse is LOW
    call     delay5us          ;wait for 5us for data stabilization
    btfsc     PORTB, KDATA
    goto      KERROR          ;if START BIT is not Zero ERROR
;START Detected
;8-bit Data Check
    movlw    0x08
    movwf    Bitcount        ;Read 8 times for 8-bit Data
RXNEXT
    bcf      STATUS, CARRY    ;Clear the Carry Bit
    rrf     DATAreg         ;rotate to the right
CKHIGH
    btfss     PORTB, CLOCK    ;Wait for CLOCK to back to High
    goto     CKHIGH
CKLOW
    btfsc     PORTB, CLOCK    ;wait for CLOCK now to LOW
    goto     CKLOW
    call     delay5us        ;5us delay
    btfsc     PORTB, KDATA    ;DATA line reading. 0 or 1
    bsf     DATAreg, MSB     ;1? Then set the MSB
    decfsz   Bitcount
    goto     RXNEXT
;Check for Parity Bit
;Wait for CLOCK back to High

```

```

CKHIGH2
    btfss    PORTB, CLOCK    ;Wait for CLOCK to back to High
    goto    CKHIGH2
CKLOW2
    btfsc    PORTB, CLOCK    ;wait for CLOCK now to LOW
    goto    CKLOW2
    call    delay5us        ;5us delay
    btfsc    PORTB, KDATA    ;Parity Bit
    goto    OneP            ;Pbit=1
    bcf     Kstat,0x00       ;Pbit=0
    goto    Stopcheck
OneP    bsf     Kstat, 0x00   ;Parity bit=1 flag
Stopcheck
;wait for CLOCK back to High
CKHIGH3
    btfss    PORTB, CLOCK    ;Wait for CLOCK to back to High
    goto    CKHIGH3
CKLOW3
    btfsc    PORTB, CLOCK    ;wait for CLOCK now to LOW
    goto    CKLOW3
    call    delay5us        ;5us delay
    btfss    PORTB, KDATA    ;STOP bit
    goto    KERROR         ;if STOP=0 , ERROR
    return
KERROR
    bsf     KSTAT, 0x02     ;ERROR FLAG set
    return

```

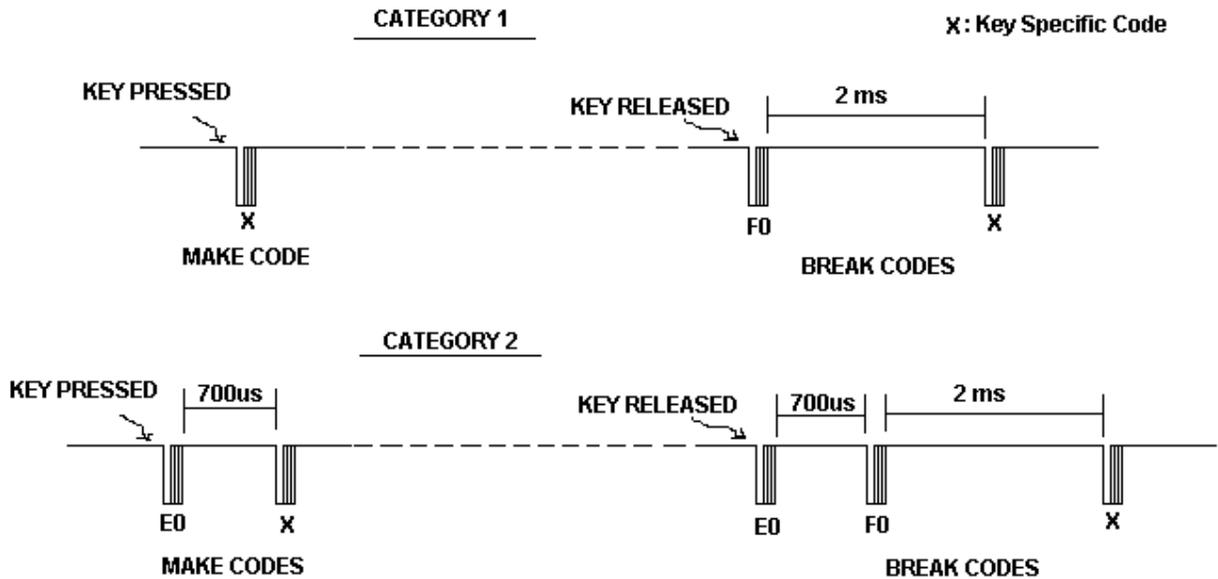
As you see from the subroutine, the code is not far from the one we developed for asynchronous subroutine for data reception. Only difference is, here, we read the data bit by monitoring the clock transition and this is the essence of the synchronous serial communication. Since we built the basic building block of a frame read, now we have to look at the Scan Codes of keyboard to know what codes are transmitted when a key is pressed and released. When a key is pressed, the keyboard controller transmits one or two 1-byte "Make" code, and when the key is released it transmits two or three 1-byte "Break" codes.

Most of the keys in a keyboard (Category 1), thus, generate 1 Make code and 2 Break codes, when they are pressed and released. The first Break code is always F0h, and the second Break code is the same as the Make code. The Make code is separated from the Break codes only by how long the key is being pressed. If that key is kept on being pressed, only the 1 Make code would be continuously generated. However, the two Break codes are separated by about 2ms.

Break code is to know when a key is actually released especially for Shift keys. While a Shift key is pressed, an 'A' would generate 'A', however, when the Shift key is released, an 'A' would be interpreted as 'a', instead.

Some keys, like HOME, DEL, Page Up, Page Down, ←, ↑, →, and ↓ (Category 2), generate two Make codes and three Break codes. For this group of keys, the first Make code is always E0. The first two Break codes are E0h and F0h, and the last break code is the same as the second Make code. Two Make codes are separated by about 700μs. The second Break code comes 700μs later after the first Break code. The last Break code arrives 2ms after the second Break

code.



The following illustration shows the key specific codes for a keyboard. Category 2 keys are shaded. There are two keys which are not in either group: Print Screen/Sys Request and Pause keys. Print Screen key has the following codes, all in hexadecimal numbers (note that key released point is indicated by a vertical bar (|)): E0, 12, E0, 7C, |, E0, F0, 7C, E0, F0, 12. The Pause key has only Make codes: E1, 14, 77, E1, F0, 14. More details on Keyboard Scan Code can found in the Microsoft's Keyboard Scan Code Specification.



Fig. 77 Key Specific Codes for a Keyboard

Therefore, according to the Scan Code illustration, the 'A' without Shift key would produce: 1C | F0, 1C. On the other hand, if we press left Shift key and holding it until we type and release 'A', and then release the left Shift key, it would generate the following codes: 12 (Left Shift Make), 1C ('A' Make), | F0 (Break), 1C ('A' Break), | F0 (Break), 12(Left Shift Break).

If you keep press the Space bar for a few seconds followed by the release, it would generate the following codes:

29 (Space Make), 29 (Space Make), 29(Space Make),|, F0 (Break), 29 (Space Break).

On the other hand, if you press the Up Arrow (\uparrow) key and release it, it would generate: E0, 75(\uparrow Make), |, E0, F0 (Break), 75 (\uparrow Break).

3. First Code - Display of the Key Code Sequence

As we just examined the code generation by keys, there are two categories considered in displaying the codes of the keys on a PC monitor. If the first code from the keyboard (the Make code) is E0, it belongs to the Category2. In category 2, we have to read at least 5 Make and Break codes. When the category key is held pressed, the Make code should be read continuously until a Break is detected. If the first Make code is not E0, it belongs to the Category 1. In Category, we have to read 3 Make and Break codes in addition to the Make codes when the key is held pressed. In Category 1, there is one exception: Left and Right Shift Keys. Usually, the Shift Make code is followed by another Make code of the Category 1 key. Therefore, we must read 5 Make and Break codes in addition to the repeated Shift Make codes and/or the repeated Make codes of the Category 1 code.

Before discussing the main part of the code, for key reading and displaying the Make and Break codes of the key, we briefly study a few subroutines that we have not discussed so far. The 5 μ s delay to read a bit after the High-to-Low transition of the CLOCK pulse, comes simply by having 10 `nop` (No Operation) instructions. Each instruction in 20 MHZ clock takes 0.2 μ s.

```
;5us Delay Subroutine
delay5us
;need total 10 insturctions
    nop
    return
```

The code display subroutine, `Kdisplay`, is to change 1-byte hex number into 2-digit ASCII codes, and transmit the ASCII codes via the serial communication module of 16F877. This subroutine is almost identical to previous Monitor display routine. Only a slight variation and variable changes has been made caused by the keyboard reading accommodation.

```
;SUBROUTINE Kdisplay
Kdisplay
```

```

    banksel    Kreg
    movwf     Kreg
    movf      Kreg,0
    movwf     Ktemp
    swapf     Ktemp,0      ;SWAP upper and lower nibbles --->W
    andlw     0x0F         ;Mask off upper nibble

    call      HTOA         ;HEX---->ASCII
    movwf     DATA1      ;First Digit

    movf      Kreg,0
    andlw     0x0F         ;mask of upper nibble

    call      HTOA         ;HEX---->ASCII
    movwf     DATA2      ;Second Digit

;TX ROUTINE FOR Display

    movf      DATA1,0
    call      Txcall       ;Serial Transmission to PC monitor
    movf      DATA2,0
    call      Txcall       ;Serial Transmission to PC monitor

    return

```

One more thing we have to know about keyboard reading is that when the keyboard is powered up it automatically sends something called BAT (Basic Assurance Test) code which basically informs the host of the status of itself. Part of the BAT is an initial set-up for the selection of Brake only, Break/Make code, and Typematic. These pieces of information are generated about 1 second after the power on. So in the code we will examine, we may want to have at least 1 second delay after the power-on reset. The follow example code is to read a keyboard and then display the key using a PC monitor, by running the `hyperterminal` at the PC side. No subroutine is listed in the code.

```

;kbd1.asm
;
;MAIN FOCUS: DISPLAYING THE KEYS PRESSED on A PC MONITOR
;
;This program is to:
;1. Read At or PS/2 Type Keyboard
;2. Display the Make and Break codes, in HEX format, on PC monitor
;
;
;   Baud rate for this is set as 19200 for Monitor display
;
;Keyboard has bi-directional synchronous serial communication
;with clock of LOW period of 30us and High period of 40us
;Data comes with 11 bits:
;1 start bit of Low
;8 data bits (LSB first)
;1 Odd parity bit
;1 Stop bit of High
;

```

```

;Data reading is done when clock goes High-to-Low transition
;
;PIC can block the data from Keyboard by pulling down the clock Low >100ms
;
;
;To leave CLOCK line High-Z, set the port Input
;To make the CLOCK line LOW, set the port as output and write 0 to the port.

;
;Algorithm
;1. Check the CLOCK pin (RB7)
;2. When CLOCK goes to LOW, read the DATA line (RB6)
;3. START>8-bit Data>1 Parity>Stop
;4. Display the result
;

;Make /Break format
;(1) (X) | (F0) (X) =====> CAT1
; L-SHIFT followed by a Key:  12 (X) |(F0) (X) |(F0) 12
; R-SHIFT followed by a key:  59 (X) |(F0) (X) |(F0) 59
;(2) (E0)(X) | (E0)(F0) (X) =====> CAT2
;
;
;Terminal set up: 8N1 19200
;
;Asynchronous mode
;
        list P = 16F877

STATUS      EQU    0x03
CARRY       EQU    0x00
ZERO        EQU    0x02
TRISB       EQU    0x86
PORTB       EQU    0x06
TXSTA       EQU    0x98      ;TX status and control
RCSTA       EQU    0x18      ;RX status and control
SPBRG       EQU    0x99      ;Baud Rate assignment
TXREG       EQU    0x19      ;USART TX Register
RCREG       EQU    0x1A      ;USART RX Register
PIR1        EQU    0x0C      ;USART RX/TX buffer status (empty or full)
RCIF        EQU    0x05      ;PIR1<5>: RX Buffer 1-Full 0-Empty
TXIF        EQU    0x04      ;PIR1<4>: TX Buffer 1-empty 0-full
TXMODE      EQU    0x20      ;TXSTA=00100000 : 8-bit, Async
RXMODE      EQU    0x90      ;RCSTA=10010000 : 8-bit, enable port, enable RX
BAUD        EQU    0x0F      ;0x0F (19200), 0x1F (9600)
CARRY       EQU    0x00
ZERO        EQU    0x02
MSB         EQU    0x07
CLOCK       EQU    0x07      ;from Keyboard
KDATA       EQU    0x06      ;from Keyboard

;
;RAM AREA

        CBLOCK      0x20
                KSTAT
                DATAreg

```

```

        MAKEreg

        BREAKreg1
        BREAKreg2
        BREAKreg

        Sbit
        Kount120us    ;Delay count (number of instr cycles for delay)
        Kount100us
        Kount1ms
        Kount10ms
        Kount100ms
        Kount1s
        Kount10s
        Kount1m
        Bitcount      ;data bit count
        Kount         ;Delay count (number of instr cycles for delay)
        DATAtemp      ;for ASCII conversion
        DATA1
        DATA2
        ASCIIreg
        Kreg
        Ktemp
ENDC

;=====
        org          0x0000
        GOTO         START
;=====
        org          0x05
;start of the program from $0005
START

        banksel     TRISB
; 1100 0000

        movlw       B'11000000' ;RB7 for CLOCK and RB6 for DATA as inputs
        movwf       TRISB
        call        ASYNC_mode
        call        delay1s          ;Give Keyboard to send STATUS to the host
;KBD initial set-up by itself
;BAT(Basic Assurance Test) code
;typematic/make/break coding

BEGIN
        banksel     TXREG
        clrf        TXREG
        banksel     DATAreg          ;the RX11bit result here
        clrf        DATAreg
        clrf        Kreg
        clrf        ASCIIreg          ;ASCII equivalent here
        clrf        Ktemp

; CHECK IF THE CLOCK is HIGH at least for 100 mS
;to make sure it does not read in the middle of data/clock stream

```

```

    banksel    PORTB
    btfss     PORTB, CLOCK
    goto      BEGIN      ;if CLOCK is LOW, start again
    call      Delay100ms ;100mS delays

;check again for CLCOK
    btfss     PORTB, CLOCK
    goto      BEGIN
;READY FOR CLOCK PULSES
;KSTAT
;KSTAT<0>: Parity Bit Value
;KSTAT<2>: Kbd error
    clrf     KSTAT
KEYIN
;X reading
    call      RX11bit          ;reading a frame

    clrf     STATUS
    movf     DATAreg,0        ;Break Code?
    xorlw    0xF0
    btfss    STATUS,ZERO
    goto     CAT
;BREAK, before MAKE code, detected. Abort It. Resume It
    goto     BEGIN
;Category detection
CAT    clrf     STATUS
    movf     DATAreg,0
    xorlw    0xE0
    btfsc    STATUS,ZERO
    goto     CAT2
    clrf     STATUS
    movf     DATAreg,0        ;L-SHIFT Key Detection
    xorlw    0x12
    btfsc    STATUS,ZERO
    goto     LRSHIFT
    clrf     STATUS
    movf     DATAreg,0        ;R-SHIFT key detection
    xorlw    0x59
    btfsc    STATUS,ZERO
    goto     LRSHIFT

                                ;L Shift ==>12 | F0 12
                                ;R Shift ==>59 | F0 59
;CAT1 has the format of (X)|(F0)(X)
CAT1   movf     DATAreg,0
    movwf    MAKEreg
    call     Kdisplay          ;(X)

;(F0) detection
    call     RX11bit
    clrf     STATUS
    movf     DATAreg,0
    xorlw    0xF0
    btfss    STATUS,ZERO
;Key is not broken. Still pressed,
    goto     CAT1            ;IF No BREAK code, Key is still pressed.
;Key is broken

```

```

        movf      DATAreg,0
        movwf    BREAKreg1
        call     Kdisplay          ;F0
;Last (X) reading
        call     RX11bit
        movf      DATAreg,0
        movwf    BREAKreg2
        call     Kdisplay          ;(X)

        call     CRLF
        call     CRLF
        goto     BEGIN             ;Read next Key

;CAT2 format (E0)(X)|(E0)(F0)(X)
CAT2   movf      DATAreg,0
        call     Kdisplay          ;E0
        call     RX11bit
        movf      DATAreg,0
        movwf    MAKEreg
        call     Kdisplay          ;(X)
; KEY still PRESSED or BROKEN
        call     RX11bit
        movf      DATAreg,0
        clrf     STATUS
        xorlw    0xE0
        btfss   STATUS,ZERO
; NOT BROKEN
        goto     CAT2
;BROKEN
        movf      DATAreg,0
        call     Kdisplay          ;E0
        call     RX11bit
        movf      DATAreg,0
        call     Kdisplay          ;F0
        call     RX11bit
        movf      DATAreg,0
        movwf    BREAKreg
        call     Kdisplay          ;(X)

        call     CRLF
        call     CRLF
        goto     BEGIN             ;Read next Key

;L-SHIFT and R-SHIFT has the form
;L-SHIFT and a character 12 X | F0 X |F0 12
;R-SHIFT and a character 59 X | F0 X |F0 59

LRSHIFT
        movf      DATAreg,0
        movwf    MAKEreg
        call     Kdisplay          ;12 or 59

;(F0) detection
        call     RX11bit
        clrf     STATUS
        movf      DATAreg,0
        xorlw    0xF0

```

```

        btfsc     STATUS,ZERO
        goto     BEGIN           ;IF no BREAK, key is not Broken yet

LRS     movf     DATAreg,0
        movwf    BREAKreg1
        call     Kdisplay       ;X

;(F0) detection
        call     RX11bit
        clrf     STATUS
        movf     DATAreg,0
        xorlw    0xF0
        btfss    STATUS,ZERO
        goto     LRS
        movf     DATAreg,0
        call     Kdisplay       ;F0
;Last (X) reading
        call     RX11bit
        movf     DATAreg,0
        movwf    BREAKreg2
        call     Kdisplay       ;(X)

        call     RX11bit
        movf     DATAreg,0
        call     Kdisplay       ;(F0)

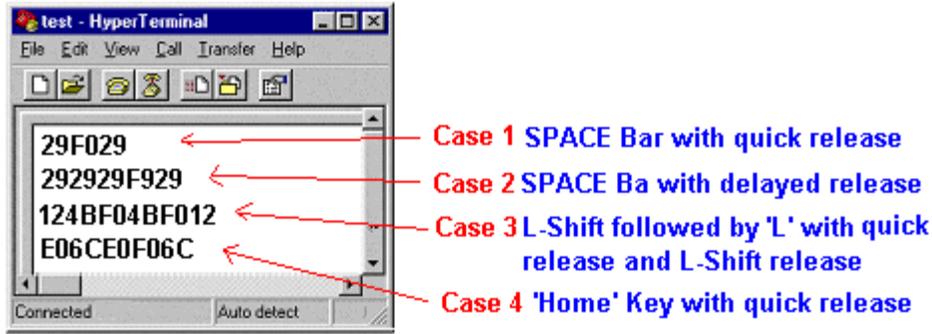
        call     RX11bit
        movf     DATAreg,0
        call     Kdisplay       ;12 or 59
        call     CRLF
        call     CRLF
        goto     BEGIN         ;Read next key

;
; SUBROUTINES
;HERE
        END
;END OF THE CODE

```

Try several keys using a keyboard and see if you get the following or similar monitor display.

1. 'Space' Bar press and quick release.
2. 'Space' Bar press and delayed release.
3. 'Shift' and 'L' followed by L release and Shift Release.
4. 'Home' key.



4. Second Code - Display of Key Itself

The second version of the program is to display the key on the monitor, not the Make and Brake codes of the key. To do this, we have to closely examine the code of a key and a combination of keys.

Category 1 Keys (Keypad is ignored)

Key	Make/Break Code (hex)	Without Shift or Caps Lock		With Shift or Caps Lock	
		Character to be Displayed	ASCII code	Character to be Displayed	ASCII code
~ `	0E	` (Apostrophe)	60	~ (Tilde)	7E
!1	16	!	31	!	21
@2	1E	@	32	@	40
#3	26	#	33	#	23
\$4	25	\$	34	\$	24
%5	2E	%	35	%	25
^6	36	^	36	^	5E
&7	3D	&	37	&	26
*8	3E	*	38	*	2A
(9	46	(39	(28
)0	45)	30)	29
-_	4E	-	2D	_	5F
+ =	55	=	3D	+	2B
BS	66	BS	08	BS	08
A	1C	a	61	A	41
B	32	b	62	B	42
C	21	c	63	C	43
D	23	d	64	D	44
E	24	e	65	E	45
F	2B	f	66	F	46
G	34	g	67	G	47
H	33	h	68	H	48
I	43	i	69	I	49
J	3B	j	6A	J	4A
K	42	k	6B	K	4B
L	4B	l	6C	L	4C
M	3A	m	6D	M	4D
N	31	n	6E	N	4E
O	44	o	6F	O	4F
P	4D	p	70	P	50

	Q	15	q	71	Q	51
	R	2D	r	72	R	52
	S	1B	s	73	S	53
	T	2C	t	74	T	54
	U	3C	u	75	U	55
	V	2A	v	76	V	56
	W	1D	w	77	W	57
	X	22	x	78	X	58
	Y	35	y	79	Y	59
	Z	1A	z	7A	Z	5A
	L-Shift	12				
	R-Shift	59				
	Enter	5A	CR	0D	CR	0D
	Space	29	Space	20	Space	20
	Cap	58				
	{ [54	[5B	{	7B
	}]	5B]	5D	}	7D
	\	5D	\	5C		7C
	: ;	4C	;	3B	:	3A
	" '	52	'	27	"	22
	< ,	41	,	2C	<	3C
	> .	49	.	2E	>	3E
	? /	4A	/	2F	?	3F

As we see from the Category 1 key table, relating its Make/Break code and matching ASCII code of the corresponding character, it is apparent that the Make/Break codes of alphanumeric keys do not match with ASCII codes of the characters of the keys. When we type a key 'A' with or without pressing a Shift key, the Make/Break code the host would get is 1Ch. This code must be changed to either 61h (ASCII code for 'a') without Shift or 41h (ASCII code for 'A') with Shift pressed. Since there is no pattern to easily convert a Make/Break code to ASCII code, we have to rely on a look-up table approach.

The approach here is to use a Make/Break code as the address where its ASCII equivalent code is stored. In other words, the Make/Break code will direct where to get the ASCII equivalent code. This sounds very simple with one minor constraint. A key in the keyboard generate the same Make/Break code, however, depending upon the Shift key or Cap Lock key, it has two ASCII equivalent codes. Therefore, we have to have two look-up tables, one without Shift (or Cap) key, and the other with Shift (or Cap) key.

The following table summarizes the two look-up tables, rearranged in the rising order of the Make/Break codes, the above table of Category 1 key.

Make/Break Code (hex)	Without Shift or Caps Lock		With Shift or Caps Lock	
	Character to be Displayed	ASCII code	Character to be Displayed	ASCII code
0E	`	60	~	7E
15	q	71	Q	51
16	!	31	!	21
1A	z	7A	Z	5A

1B	s	73	S	53
1C	a	61	A	41
1D	w	77	W	57
1E	2	32	@	40
21	c	63	C	43
22	x	78	X	58
23	d	64	D	44
24	e	65	E	45
25	4	34	\$	24
26	3	33	#	23
29	Space	20	Space	20
2A	v	76	V	56
2B	f	66	F	46
2C	t	74	T	54
2D	r	72	R	52
2E	5	35	%	25
31	n	6E	N	4E
32	b	62	B	42
33	h	68	H	48
34	g	67	G	47
35	y	79	Y	59
36	6	36	^	5E
3A	m	6D	M	4D
3B	j	6A	J	4A
3C	u	75	U	55
3D	7	37	&	26
3E	8	38	*	2A
41	,	2C	<	3C
42	k	6B	K	4B
43	i	69	I	49
44	o	6F	O	4F
45	0	30)	29
46	9	39	(28
49	.	2E	>	3E
4A	/	2F	?	3F
4B	l	6C	L	4C
4C	;	3B	:	3A
4D	p	70	P	50
4E	-	2D	=	5F
52	'	27	"	22
54	[5B	{	7B
55	=	3D	+	2B
5A	CR	0D	CR	0D
5B]	5D	}	7D
5D	\	5C		7C
66	BS	08	BS	0B

Then, how do we generate a table or two tables in 16F877 assembly language programming environment? The easiest way of table building is to change the program counter (PC). As we all know, PC indicates the next address to fetch a program code and execute. PC in 16F877 is a 13-bit register which has address access range of $2^{13}=8K$ words. The lower 8 bits of the PC is can be controlled (i.e., read from and written to) by PCL (PC Lower Byte) register. The upper 5

bits are not directly accessed by PCH (PC High Byte), instead it is controlled by PCLATH (PC Latch) register.

Table formation is utilized by the PCL in a subroutine, along with an instruction (`retlw k`, return with `k` in `W` register) or a Microchip Assembler (MPASM) directive (`DT`, Define Table). In other words, a table formation is a subroutine building with PCL, and lines of `retlw` and/or `DT`.

Let's have a simple program to illustrate how to form a table. In the previous example of serial communication, we have had a key typed in the keyboard echoed on a PC monitor. Now we want to change the program slightly so that the program receives only numbers (0 to 9) from the keyboard and echoes corresponding characters determined by the following table:

Key	0	1	2	3	4	5	6	7	8	9
Echo	A	B	C	D	E	F	G	H	I	J

The pseudo-code of the revised program would go like this:

- (a) Receive a number from keyboard by calling `RXPOLL` subroutine
- (b) The received key is converted from ASCII (30h for '0', for example) to a hex number (00h for '0', for example).
- (c) Check the table and find the corresponding character ('A' for 00h for example).
- (d) Transmit the matched character to PC monitor by calling `TXPOLL` subroutine.

The following code lists the table subroutine, `Keytable`:

```

;
Keytable
    addwf    PCL
    retlw   'A'    ;PC+0
    retlw   'B'    ;PC+1
    retlw   'C'    ;PC+2
    retlw   'D'
    retlw   'E'
    retlw   'F'
    retlw   'G'
    retlw   'H'
    retlw   'I'
    retlw   'J'          ;PC+9

```

When this subroutine is called, the return address is stored in the Stack, and the content of PC is the starting address of this subroutine in other words, the PC is the address of the first instruction line is the subroutine:

```
addwf PCL.
```

That's why when a subroutine is called, it's executed from the first line of the subroutine. When the `addwf PCL` is executed (sum of the content of `W` and `PCL`), the PC changes by the amount of `W` content. If the `W` is 0, then there is no change in PC so the next line

```
retlw 0x41
```

would be executed. (Remember PC always directs the next code to fetch and execute.)

The instruction

```
retlw k
```

is combination of two instructions:

```
movlw k
return
```

Therefore, `retlw 0x41` returns to the next line after the caller with 41h in the W register. If we place this content in the TXREG of USART module and call TXPOLL, the character 'A' would be displayed on the monitor.

What happens if number 5 is pressed in the keyboard? The hex-converted number 05h would be placed in W register after RXPOLL subroutine, then when the `Keytable` is called, PCL is added by 5, therefore, the sixth line after the PCL instruction would be executed, which results in the display of 'F' on the monitor. The following code lists the main part and the table of the whole program.

```

;=====
    org      0x0000
    GOTO    START
;=====
    org      0x05
;start of the program from $0005
START
    call    ASYNC_mode      ;initialization of USART module

BEGIN
    banksel TXREG
    clrf   TXREG
    clrf   RCREG
    clrf   Ktemp

AGAIN
    call   RXPOLL          ;read a key
    movwf Ktemp
    movlw 0x30
    subwf Ktemp,0         ;Make it into a hex number (f-W -->d)
    call   Keytable       ;Call Keytable match
    call   TXPOLL         ;matched character
    call   CRLF           ;CR and LF
    goto  AGAIN          ;Repeat

Keytable
    addwf  PCL
    retlw  'A'           ;PC+0
    retlw  'B'           ;PC+1
    retlw  'C'           ;PC+2
    retlw  'D'
    retlw  'E'
    retlw  'F'
    retlw  'G'

```

```

retlw    'H'
retlw    'I'
retlw    'J'

```

The keytable can also be formed by using an assembler directive, DT. DT actually is a multiple of retlw instructions:

```

Keytable
    addwf    PCL
    DT      "ABCDEFGHJIJ"    ;PC+0, 1, 2, 3, 4, 5, 6, 7, 8, and 9

```

There is one important matter to remember when we use the PCL writing instruction, addwf PCL. When a PCL is written, the lower 5 bits of PCLATH register is also written to PCH. Since PC is PCH (PC[12:8]) concatenated by PCL (PC[7:0]), if both the location of the caller line and that of the keytable subroutine are located inside the 00h – FFh boundary, there is no problem, since PCH for both parts is the same. Let's discuss this using the List file of the above program.

```

MPASM 02.61 Released          19TABLE1.ASM    5-20-2004  14:58:05
LOC  OBJECT CODE          LINE SOURCE TEXT

                                00036
;=====
0000          00037          ORG      0X0000
0000  2805          00038          GOTO   START
                                00039
;=====
0005          00040          ORG      0X05
                                00041 ;start of the program from $0005
0005          00042          START
0005  201E          00043          CALL   ASYNC_MODE
                                00044
                                00045
0006          00046          BEGIN
0006  1283 1303          00047          BANKSEL TXREG
0008  0199          00048          CLRF   TXREG
0009  019A          00049          CLRF   RCREG
000A  01A0          00050          CLRF   KTEMP
000B          00051          AGAIN
000B  2033          00052          CALL   RXPOLL
000C  00A0          00053          MOVWF  KTEMP
000D  3030          00054          MOVLW  0X30
000E  0220          00055          SUBWF  KTEMP,0
000F  2013          00056          CALL   KEYTABLE
0010  202B          00057          CALL   TXPOLL
0011  203B          00058          CALL   CRLF
0012  280B          00059          GOTO   AGAIN
                                00060
                                00061
0013          00062          KEYTABLE
0013  0782          00063          ADDWF  PCL
0014  3441          00064          RETLW  'A'    ;PC+0
0015  3442          00065          RETLW  'B'    ;PC+1
0016  3443          00066          RETLW  'C'    ;PC+2

```

0017	3444	00067	RETLW	'D'
0018	3445	00068	RETLW	'E'
0019	3446	00069	RETLW	'F'
001A	3447	00070	RETLW	'G'
001B	3448	00071	RETLW	'H'
001C	3449	00072	RETLW	'I'
001D	344A	00073	RETLW	'J'

The first column is the program memory address and the second and the thirds are for the op-codes of the instruction. As we see from the List file, the caller (which calls the `Keytable` subroutine) is located at 000F. In other words, up until the `Keytable` is called, the PC is, in binary number, 0 0000 0000 1111 (or 000Fh). In other words, the upper byte of PC, PCH, is 00000 (or 00h). This would be the content of the lower 5 bits of PCLATH when the `Keytable` is called. When the PCL writing is done, 00h, the content of PCLATH will be filled to PCH, making the upper byte of PC zero. This, however, does not cause any trouble, because the `Keytable` subroutine, from 0013h to 001D, is with only lower byte portion of PC, PCL. In other words, the PC access to the subroutine instructions are inside the range with PCH=00h.

If the caller and the `keytable` subroutine is apart more than FFh each other, and there is uncertainty about this, we have to see the List file and decide what action should be made to avoid possible PC related problem.

For example, consider that the caller is located in 000Fh (with the current PC) and the `keytable` subroutine is located in 0113h instead. When the subroutine is called, the content of PC is 0000Fh. Therefore the PCH is 00000b. This binary value of 00000h would be written to the PCH portion of the PC when PCL writing is performed. Therefore the final value of PC would be 0013h, instead of the desired value of 0113h. Therefore, we have to have the following instruction just above the PCL writing instruction:

```
bsf   PCLATH, 0x00
```

Check with Microchip's 16F877 manual for detailed description of PC, PCH, PCL, PCLATH, and PCL writing with relation to the four different pages of 8K word program memory structure.

Now, we are ready to explore the second version of the keyboard connection to a PC monitor. The second example is to display the characters of the keys itself not the Make/Break codes of the keys. We use the Make/Break code table for `keytable` formulation. Since we have to consider the three following conditions; (a) when Shift key is not pressed, (b) when Shift key is pressed, and (c) when Caps Lock key is pressed. When Caps Lock is pressed we change only alphabets to upper cases and keep all other keys as if no Shift key is pressed.

The first table is for the keys without a Shift key, `NoShiftKeyTable`. According to the Make/Break code, the lowest hex number of the code is 0E which, without a Shift Key, is for '' (Apostrophe), with its ASCII code 60. Then there are gaps until we have the next Make/Break code, which are 15h for 'q' and 16h for 'l'. The following code shows the table for no Shift key. We will start the table 0100h. Since we assume that the caller is located before 0100h address, before the call is made, the PCH would be 00h. At the first line of the subroutine, we configure the PCH of PC by PCLATH instruction.

The second table is for the keys with a Shift key, `ShiftKeyTable`, located just after the first table. This table can be similarly formed.

The third table accommodates Caps Lock key, `CapKeyTable`, which starts from 0200h. This address (also PC) requires additional PCL configuration. Details are commented in the code section.

The pseudo-code of the main section of the program goes like this. It's rather complicated. So we have to fully understand the Make/Break code pattern when a key is pressed and released.

- (1)BEGIN: read the first frame and get the 8-bit key data
- (2)Check if the key is 12h (Left Shift) or 59h (Right Shift). If it is, jump to LRSHIFT
- (3) Check if the key is 58h(Caps Lock). If it is, jump to CAPS
- (4)NOSHIFT:
 - (a) Call `NoShiftKeyTable` and display what W register holds
 - (b) read the next frame. If the key is F0h, then read one more frame then go to BEGIN
 - (c) If the key is not F0h, go to NOSHIFT
- (5) LRSHIFT:
 - (a)Read the next frame. If the key is F0h, then read one more frame then go to BEGIN
 - (b)If the key is 12 or 59h (Shift key is not broken yet), go to LRSHIFT
 - (c) If the key is neither 12h nor 59h, call `ShiftKeyTable` and display what W register holds.
 - (d)read the next frame. If the key is not F0h, call `ShiftKeyTable` and display
 - (e) if the key is F0h, go to LRS
 - (f)LRS: read next frame, if the key is 12h, go to BEGIN. If the key is not 12h, go to LRSHIFT.
- (6)CAPS:
 - (a)Read the next two frames.
 - (b) CAPNEXT: Read the next frame. If the key is 58h, then read two more frames, then go to BEGIN
 - (b)If the key is not 58h, call `CapKeyTable` and display what W register holds.
 - (c)Read the next frame. If the key is F0h, read the next frame and go to CAPNEXT.
 - (d) If the key is 58h, then read two more frames, then go to BEGIN.

The next listing shows a complete code for displaying characters, numbers, and other symbols while accommodating Shift (Left and Right) and Caps Lock keys. Follow each line and comment closely to better understand the program.

```

;kbd3.asm
;
;SHIFT and CAPS LOCK are featured
;
;This program is to:
;1. Read At or PS/2 Type Keyboard
;2. Display them, as characters, on PC monitor
;
;3. Note that all Category 2 keys (E0 keys) are ignored

```

```

;
;   Baud rate for this is set as 19200 for Monitor display
;
;
;Terminal set up: 8N1 19200
;
;Asynchronous mode
;

        list P = 16F877
PCL      EQU    0x02      ;For Key Table Calling (Lower PC)
PCLATH   EQU    0x0A      ;For upper part of PC
STATUS   EQU    0x03
CARRY    EQU    0x00
ZERO     EQU    0x02
TRISB    EQU    0x86
PORTB    EQU    0x06
TXSTA    EQU    0x98      ;TX status and control
RCSTA    EQU    0x18      ;RX status and control
SPBRG    EQU    0x99      ;Baud Rate assignment
TXREG    EQU    0x19      ;USART TX Register
RCREG    EQU    0x1A      ;USART RX Register
PIR1     EQU    0x0C      ;USART RX/TX buffer status (empty or full)
RCIF     EQU    0x05      ;PIR1<5>: RX Buffer 1-Full 0-Empty
TXIF     EQU    0x04      ;PIR1<4>: TX Buffer 1-empty 0-full
TXMODE   EQU    0x20      ;TXSTA=00100000 : 8-bit, Async
RXMODE   EQU    0x90      ;RCSTA=10010000 : 8-bit, enable port, enable RX
BAUD     EQU    0x0F      ;0x0F (19200), 0x1F (9600)
CARRY    EQU    0x00
ZERO     EQU    0x02
MSB      EQU    0x07
CLOCK    EQU    0x07      ;from Keyboard
KDATA    EQU    0x06      ;from Keyboard

;
;RAM AREA

        CBLOCK    0x20
                TXtemp
                KSTAT
                DATAreg
                Kount100us
                Kount10ms
                Kount1s
                Bitcount      ;data bit count
        ENDC

;=====
        org      0x0000
        GOTO    START
;=====
        org      0x05
;start of the program from $0005
START

```

```

        banksel    TRISB
; 1100 0000

        movlw     B'11000000' ;RB7 for CLOCK and RB6 for DATA as inputs
        movwf    TRISB
        call     ASYNC_mode
        call     delays      ;Give Keyboard to send STATUS to the host
;KBD initial set-up by itself
;BAT(Basic Assurance Test) code
;typematic/make/break coding

BEGIN
        banksel    TXREG
        clrf      TXREG
        banksel    DATAreg
        clrf      DATAreg

; CHECK IF THE CLOCK is HIGH at least for 10mS

        banksel    PORTB
        btfss     PORTB, CLOCK
        goto      BEGIN      ;if CLOCK is LOW, start again
        call     Delay10ms   ;10mS delays
;This short delay speeds up the response
;check again for CLCOK
        btfss     PORTB, CLOCK
        goto      BEGIN
;READY FOR CLOCK PULSES
;KSTAT
;KSTAT<0>: Parity Bit Value
;KSTAT<2>: Kbd error
        clrf      KSTAT
KEYIN
;X reading
        call     RX11bit      ;
        clrf      STATUS
        movf     DATAreg,0   ;Break Code?
        xorlw    0xF0
        btfss   STATUS,ZERO
        goto     CAT
;BREAK is detected. Abort It. Resume It
        goto     BEGIN
;Category detection (SHIFT or CAPS LOCK)
CAT    clrf      STATUS
        movf     DATAreg,0
        xorlw    0xE0
        btfsc   STATUS,ZERO
        goto     Begin      ;E0 keys (CAT2) are ignored
        clrf      STATUS
        movf     DATAreg,0
        xorlw    0x12
        btfsc   STATUS,ZERO
        goto     LRSHIFT
        clrf      STATUS
        movf     DATAreg,0
        xorlw    0x59

```

```

    btfsc     STATUS,ZERO
    goto     LRSHIFT
    clrf     STATUS
    movf     DATAreg,0
    xorlw    0x58          ;CAPS LOCK
    btfsc     STATUS,ZERO
    goto     CAPS

                                ;L Shift ==>12 | F0 12
                                ;R Shift ==>59 | F0 59
;CAT1 without Shift or Caps Lock Key
CAT1 movf     DATAreg,0
    call     NoShiftKeyTable      ;(X) display
    call     TXCALL
;(F0) detection
    call     RX11bit
    clrf     STATUS
    movf     DATAreg,0
    xorlw    0xF0
    btfss    STATUS,ZERO
;Key is not broken. Still pressed,
    goto     CAT1
;Key is broken
;Last (X) reading
    call     RX11bit              ;(X) after F0

    goto     BEGIN

;L-SHIFT and R-SHIFT has the form
;L-SHIFT and a character 12 X | F0 X |F0 12
;R-SHIFT and a character 59 X | F0 X |F0 59

LRSHIFT                                ;12 or 59 entered

;(F0) detection
    call     RX11bit
    clrf     STATUS
    movf     DATAreg,0
    xorlw    0xF0
    btfsc    STATUS,ZERO
    goto     BEGIN

                                ;X

    clrf     STATUS                ;if (12) do not display
    movf     DATAreg,0
    xorlw    0x12
    btfsc    STATUS, ZERO
    goto     LRSHIFT

    clrf     STATUS                ;if (59) do not display
    movf     DATAreg,0
    xorlw    0x59
    btfsc    STATUS, ZERO
    goto     LRSHIFT

                                ;a Key (X) is entered

```

```

    movf    DATAreg,0
    call    ShiftKeyTable
    call    TXCALL

;(F0) detection
    call    RX11bit
    clrf    STATUS
    movf    DATAreg,0
    xorlw   0xF0
    btfss   STATUS,ZERO
    goto    LRSHIFT
;Last (X) reading
    call    RX11bit
    movf    DATAreg,0
    clrf    STATUS           ;check if (X) or (12) entered after F0
    xorlw   0x12
    btfsc   STATUS,ZERO
    goto    BEGIN
    goto    LRSHIFT
;
CAPS           ;caps lock (58) entered

;(F0) detection
    call    RX11bit           ;this must be F0
    call    RX11bit           ;this must be (58) again

CAPnext
    call    RX11bit           ;Check if (58) or other
    clrf    STATUS
    movf    DATAreg,0
    xorlw   0x58
    btfss   STATUS,ZERO
    goto    CAPtwo           ;End of CAP session
    call    RX11bit           ;F0
    call    RX11bit           ;(58)
    goto    BEGIN
                                ;a Key (X) is entered

CAPtwo
    movf    DATAreg,0
    call    CAPKeyTable
    call    TXCALL

;(F0) detection
    call    RX11bit           ;
    clrf    STATUS
    movf    DATAreg,0
    xorlw   0xF0
    btfss   STATUS,ZERO
    goto    CAPtwo

;Last (X) reading           ;F0 is read
    call    RX11bit           ;(X) again and ignore
    goto    CAPnext

```

```

;SUBROUTINE RX11bit =====
;RX Routine for 11-bit read
;1 Start
;8 Data (LSB first)
;1 Parity (Odd)
;1 Stop (HIGH)
;KSTAT Bit Info
; KSTAT<0> : parity
; KSTAT<2>:kBD Error
RX11bit
    clrf        DATAreg
    banksel    PORTB
;Let it have at least 500us CLOCK high period
    btfss     PORTB, CLOCK
    goto      RX11bit        ;if CLOCK is LOW, start again
    call      Delay100us    ;200uS delays
    call      Delay100us
;check again for CLCOK
    btfss     PORTB, CLOCK
    goto      RX11bit
;Clock Check
Scheck
    btfsc     PORTB,CLOCK
    goto      Scheck
    call      delay5us      ;wait for 5us      for data stabilization
    btfsc     PORTB, KDATA
    goto      KERROR        ;if START BIT is not Zero ERROR
;START Detected
;8-bit Data Check
    movlw    0x08
    movwf    Bitcount      ;8 data bits
RXNEXT
    bcf      STATUS, CARRY    ;Clear the Carry Bit
    rrf      DATAreg        ;rotate to the right
CKHIGH
    btfss    PORTB, CLOCK    ;Wait for CLOCK to back to High
    goto    CKHIGH
CKLOW
    btfsc    PORTB, CLOCK    ;wait for CLOCK now to LOW
    goto    CKLOW
    call    delay5us        ;5us delay
    btfsc    PORTB, KDATA    ;0 or 1
    bsf     DATAreg, MSB    ;1? Then set the MSB
    decfsz  Bitcount
    goto    RXNEXT
;Check for Parity Bit
;Wait for CLOCK back to High
CKHIGH2
    btfss    PORTB, CLOCK    ;Wait for CLOCK to back to High
    goto    CKHIGH2
CKLOW2
    btfsc    PORTB, CLOCK    ;wait for CLOCK now to LOW
    goto    CKLOW2
    call    delay5us        ;5us delay
    btfsc    PORTB, KDATA    ;Parity Bit
    goto    OneP            ;Pbit=1
    bcf     Kstat,0x00      ;Pbit=0
    goto    Stopcheck

```

```

Onep  bsf          Kstat, 0x00 ;Pbit=1
Stopcheck
;wait for CLOCK back to High
CKHIGH3
    btfss         PORTB, CLOCK      ;Wait for CLOCK to back to High
    goto         CKHIGH3
CKLOW3
    btfsc         PORTB, CLOCK      ;wait for CLOCK now to LOW
    goto         CKLOW3
    call         delay5us          ;5us delay
    btfss         PORTB, KDATA      ;STOP bit
    goto         KERROR           ;if STOP=0 , ERROR
    return
KERROR
    bsf          KSTAT, 0x02
    return
;=====
;RX TX Initialization with Async Mode
;Async_mode Subroutine
Async_mode
    banksel     SPBRG
    movlw      baud          ;B'00001111' (19200)
    movwf     SPBRG
    banksel     TXSTA
    movlw      TXMODE        ;B'00100000' Async Mode
    movwf     TXSTA
    banksel     RCSTA
    movlw      RXMODE        ;B'10010000' Enable Port
    movwf     RCSTA
    return
;=====
TXCALL
;slight change so that CR make CR and LF together
    banksel     TXtemp
    movwf     TXtemp
    clrf      STATUS
    movf      TXtemp,0
    xorlw     0x0D
    btfsc     STATUS,ZERO
    goto     CRNLF
Keymain
    banksel     PIR1
    btfss     PIR1, TXIF    ; Check if TX buffer is empty
    goto     Keymain
    banksel     TXREG
    movf      Txtemp,0
    movwf     TXREG        ; Place the character to TX buffer
    return
CRNLF call     CRLF
    return
;===
CRLF
    banksel     PIR1
    btfss     PIR1, TXIF
    goto     CRLF
    banksel     TXREG
    movlw     H'0d'        ;CR

```

```

        movwf      TXREG
LFkey
        banksel   PIR1
        btfss    PIR1, TXIF
        goto     LFkey
        banksel   TXREG
        movlw    H'0A'      ;LF
        movwf    TXREG
        return

;=====
delay5us
;need total 10 instructions
        nop
        return

;=====

Delay100us
        banksel   Kount100us
        movlw    H'A4'
        movwf    Kount100us
R100us
        decfsz   Kount100us
        goto     R100us
        return

;
;10ms delay
; call 100 times of 100 us delay (with some time discrepancy)
Delay10ms
        banksel   Kount10ms
        movlw    H'64' ;100
        movwf    Kount10ms
R10ms  call     delay100us
        decfsz   Kount10ms
        goto     R10ms
        return

;

;1 sec delay
;call 100 times of 10ms delay
Delay1s
        banksel   Kount1s
        movlw    0x6A
        movwf    Kount1s
R1s    call     Delay10ms
        decfsz   Kount1s
        goto     R1s
        return

```



```

    retlw    0            ;+37
    retlw    0            ;+38
    retlw    0            ;+39
    DT      +"mju78"      ;+3A, 3B, 3C, 3D, 3E
    retlw    0            ;+3F
    retlw    0            ;+40
    DT      ",kio09"      ;+41, 42,43,44,45,46
    retlw    0            ;+47
    retlw    0            ;+48
    DT      ". /!;p-"      ;+49, 4A, 4B, 4C, 4D, 4E
    retlw    0            ;+4F
    retlw    0            ;+50
    retlw    0            ;+51
    retlw    0x27         ;+52  single quote
    retlw    0            ;+53
    DT      "[="          ;+54, 55
    retlw    0            ;+56
    retlw    0            ;+57
    retlw    0            ;+58
    retlw    0            ;+59
    retlw    0x0D         ;+5A  Return
    retlw    ']'          ;+5B
    retlw    0            ;+5C
    retlw    0x5C         ;+5D  \
    retlw    0            ;+5E
    retlw    0            ;+5F
    retlw    0            ;+60
    retlw    0            ;+61
    retlw    0            ;+62
    retlw    0            ;+63
    retlw    0            ;+64
    retlw    0            ;+65
    retlw    0x08         ;+66  Backspace
;end if NoShiftKetTable

;With Shift Key Table
shiftKeyTable
    bsf      PCLATH, 0x00
    bcf      PCLATH,0x01
    addwf    PCL

;Note that writing to PCL also brings the content of lower 5 bits of PCLATH
;to PC.
;In this code, this Table starts from 0134h
;While the main part of program which calls this Table is somewhere in
;0045h.
;for 0045, the PCLATH part is 00
;Therefore we have to manually set the PCLATH part so that it
;can point inside this table
    retlw    0            ;PC+0
    retlw    0            ;PC+1
    retlw    0            ;+2
    retlw    0
    retlw    0
    retlw    0
    retlw    0
    retlw    0
    retlw    0
    retlw    0

```

```

retlw      0
retlw      0
retlw      0
retlw      0
retlw      0           ;+0D
retlw      0x7E       ;+0E MAKE/BREAK= 0E ----->ASCII 7E (~)
retlw      0           ;+0F
retlw      0
retlw      0
retlw      0
retlw      0           ;+13
retlw      0           ;+14
DT         "Q!"       ;+15, 16
retlw      0           ;+17
retlw      0
retlw      0
DT         "ZSAW@"    ;+1A, 1B, 1C, 1D, 1E
retlw      0           ;+1F
retlw      0           ;+20
DT         "CXDE$#"   ;+21, 22, 23, 24, 25, 26
retlw      0           ;+27
retlw      0           ;+28
retlw      ' '        ;+29 Space
DT         "VFTR%"    ;+2A, 2B, 2C, 2D, 2E
retlw      0           ;+2F
retlw      0           ;+30
DT         "NBHGY^"   ;+31, 32, 33, 34,35,36
retlw      0           ;+37
retlw      0           ;+38
retlw      0           ;+39
DT         "MJU&*"    ;+3A, 3B, 3C, 3D, 3E
retlw      0           ;+3F
retlw      0           ;+40
DT         "<KIO)("    ;+41, 42,43,44,45,46
retlw      0           ;+47
retlw      0           ;+48
DT         ">?L:P_"    ;+49, 4A, 4B, 4C, 4D, 4E
retlw      0           ;+4F
retlw      0           ;+50
retlw      0           ;+51
retlw      0x22       ;+52 double quote
retlw      0           ;+53
DT         "{+"       ;+54, 55
retlw      0           ;+56
retlw      0           ;+57
retlw      0           ;+58
retlw      0           ;+59
retlw      0x0D       ;+5A Return
retlw      '}'        ;+5B
retlw      0           ;+5C
retlw      0x7C       ;+5D |
retlw      0           ;+5E
retlw      0           ;+5F
retlw      0           ;+60
retlw      0           ;+61
retlw      0           ;+62
retlw      0           ;+63

```

```

    retlw    0            ;+64
    retlw    0            ;+65
    retlw    0x08        ;+66 Backspace

;CAPs Lock Key Table
    org      0x0200      ;This starts from 0200h
                                ;to well arrange the PCH of PC
CAPKeyTable
    bsf      PCLATH, 0x01
    bcf      PCLATH, 0x00 ;see here that PCH is 02h?
    addwf   PCL

    retlw    0            ;PC+0
    retlw    0            ;PC+1
    retlw    0            ;+2
    retlw    0
    retlw    0x60        ;+0D
                                ;+0E MAKE/BREAK= 0E ---->ASCII = 0x60
Apostrophe
    retlw    0            ;+0F
    retlw    0
    retlw    0
    retlw    0
    retlw    0            ;+13
    retlw    0            ;+14
    DT      "Q1"          ;+15, 16
    retlw    0            ;+17
    retlw    0
    retlw    0
    DT      "ZSAW2"       ;+1A, 1B, 1C, 1D, 1E
    retlw    0            ;+1F
    retlw    0            ;+20
    DT      "CXDE43"      ;+21, 22, 23, 24, 25, 26
    retlw    0            ;+27
    retlw    0            ;+28
    retlw    ' '          ;+29 Space
    DT      "VFTR5"       ;+2A, 2B, 2C, 2D, 2E
    retlw    0            ;+2F
    retlw    0            ;+30
    DT      "NBHGY6"      ;+31, 32, 33, 34, 35, 36
    retlw    0            ;+37
    retlw    0            ;+38
    retlw    0            ;+39
    DT      "MJU78"       ;+3A, 3B, 3C, 3D, 3E
    retlw    0            ;+3F
    retlw    0            ;+40
    DT      ",KIO09"      ;+41, 42, 43, 44, 45, 46
    retlw    0            ;+47

```

```

retlw      0           ;+48
DT         ". /L;P-"  ;+49, 4A, 4B, 4C, 4D, 4E
retlw      0           ;+4F
retlw      0           ;+50
retlw      0           ;+51
retlw      0x27        ;+52  single quote
retlw      0           ;+53
DT         "[="       ;+54, 55
retlw      0           ;+56
retlw      0           ;+57
retlw      0           ;+58
retlw      0           ;+59
retlw      0x0D        ;+5A  Return
retlw      ']'         ;+5B
retlw      0           ;+5C
retlw      0x5C        ;+5D  \
retlw      0           ;+5E
retlw      0           ;+5F
retlw      0           ;+60
retlw      0           ;+61
retlw      0           ;+62
retlw      0           ;+63
retlw      0           ;+64
retlw      0           ;+65
retlw      0x08        ;+66  Backspace

;END OF CODE
      END

```

Compile and run the above example code, see how fast or slow the 16F988 responds and displays the keys.

5. Third Code - Display Key in LCD

The next, final version is to change the display medium from the PC monitor to the 20x4 LCD module, we studied in the Serial Communication. As in the digital clock, we will apply the 4-bit interface configuration to display characters.

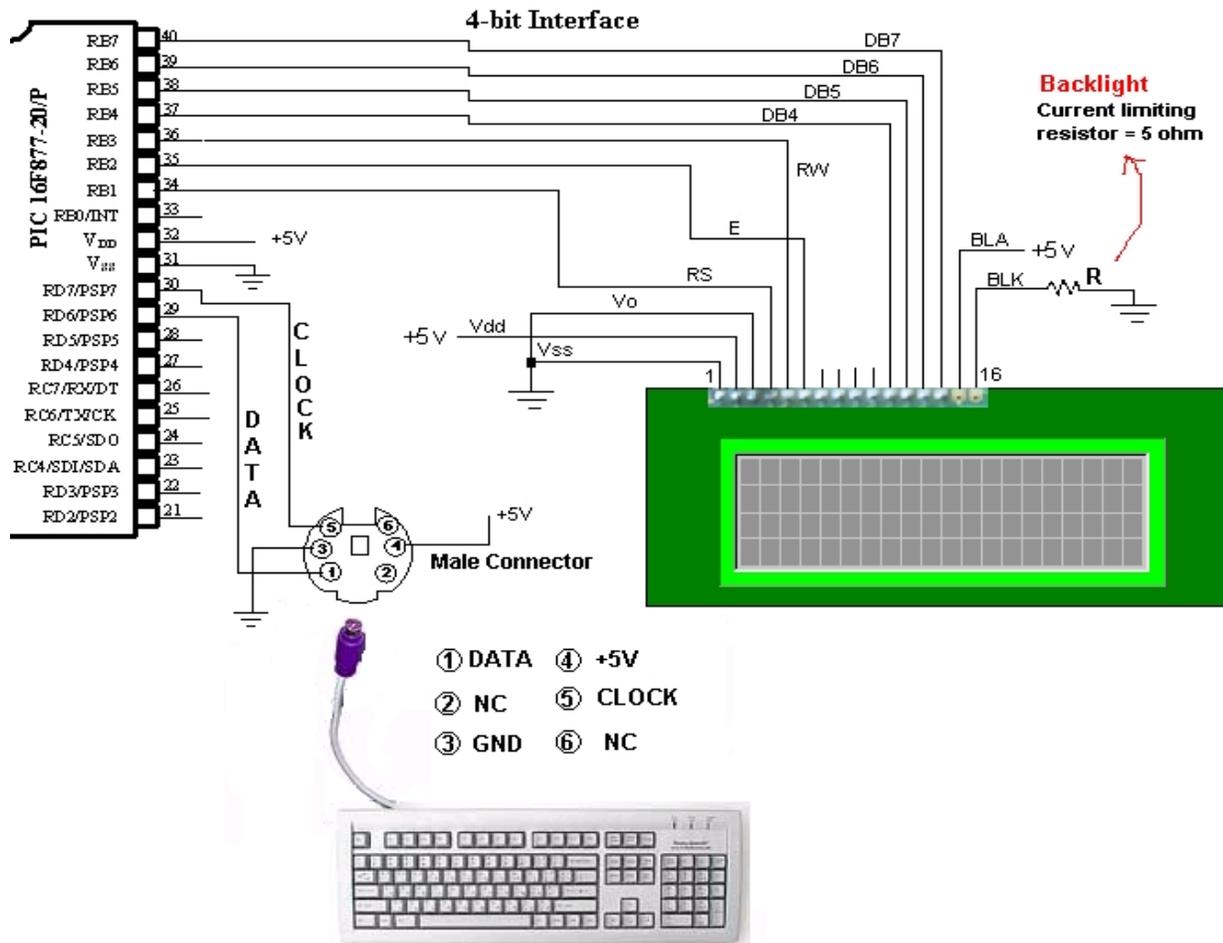


Fig. 78 Connection to display key in LCD

As shown in the schematic, PORTB is assigned to LCD control as we did before, and RD7 and RD6 are assigned to the DATA and CLOCK signal lines of the keyboard. In this final version, we will have two example codes. The first one is to display the keyboard on the LCD from the first column of the line 1 to the last column of the line 4. The first one, for convenience and simplicity, ignores Back Space (BS) and Carriage Return (CR) keys. The LCD controller inside the module does not have the stored character for BS and CR, therefore, no output will be displayed. The accommodation of these two keys is made in the second example code of this application.

Since most of the subjects here are related to the LCD control and keyboard reading, the only important thing is to remember the cursor location and its address and control them for display. In most LCD module, all the characters in ASCII code table and some other special characters are stored at the addresses which are the ASCII codes themselves. For example, the character 'A'

in dot matrix form is stored at 41h, and the ASCII code of 'A' is 41h. This means that we can use the two key tables we used for PC monitor display without any change.

As we discussed in the LCD module in Chapter 6, there is somewhat weird address allocation of 20x4 positions of the LCD module. Here we show again the address of each display cell of 20x4 LCD module:

Position	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20
First line	00h	01h	02h	03h	04h	05h	06h	07h	08h	09h	0Ah	0Bh	0Ch	0Dh	0Eh	0Fh	10h	11h	12h	13h
Second line	40h	41h	42h	43h	44h	45h	46h	47h	48h	49h	4Ah	4Bh	4Ch	4Dh	4Eh	4Fh	50h	51h	52h	53h
Third line	14h	15h	16h	17h	18h	19h	1Ah	1Bh	1Ch	1Dh	1Eh	1Fh	20h	21h	22h	23h	24h	25h	26h	27h
Fourth Line	54h	55h	56h	57h	58h	59h	5Ah	5Bh	5Ch	5Dh	5Eh	5Fh	60h	61h	62h	63h	64h	65h	66h	67h

As you see the addresses are continuous from line 1 to line 3, and from line 2 and 4. Therefore, displaying characters continuously from the first line to the last involves tracking the current cursor position and its address. For example, if the current cursor position is the 20th position of line 2 (address = 53h), the next cursor position must be the 1st position of line 3 (address = 14h).

Indeed, there is a way to read the current cursor address by reading the address from the LCD module. However, actually, reading the address after every writing a data into LCD is not necessary. In the reset, the LCD is usually configured to start from the first position of the first line, and as a character is displayed the address incremented automatically. Therefore, if we assign the cursor position, not the cursor address which must be read from LCD module, and increase the cursor position every time we write a data to the LCD module, we can easily track the current cursor position. So in the first example code for LCD display, we assign a file register CURSOR to track the current cursor position. The CURSOR's value of the cursor position, in the continuous order, is assigned as follows.

Position	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20
First line	01	02	03	04	05	06	07	08	09	0A	0B	0C	0D	0E	0F	10	11	12	13	14
Second line	15	16	17	18	19	1A	1B	1C	1D	1E	1F	20	21	22	23	24	25	26	27	28
Third line	29	2A	2B	2C	2D	2E	2F	30	31	32	33	34	35	36	37	38	39	3A	3B	3C
Fourth Line	3D	3E	3F	40	41	42	43	44	45	46	47	48	49	4A	4B	4C	4D	4E	4F	50

So when it starts, the CURSOR=1 while the cursor address is 00h. Similarly, CURSOR=3B while the cursor address in the LCD module is 26h. Also, if CURSOR=15, we change the cursor address to 40h so that the LCD module actually moves the cursor to the first position of the second line. The cursor is not changed accordingly unless the cursor address is changed according to the address table: CURSOR itself cannot change the cursor position; it is only for our convenience in cursor position tracking.

Since our LCD configuration automatically increases the cursor address by one whenever a character is displayed, we increase the CURSOR by 1 every time we write a character to the LCD module, and we check in which line the cursor is currently pointed. If the current cursor position is, for example, at the last position of line 1, then, the next CURSOR value must be changed to the first position of line 2. This seems quite trivial since CURSOR value is well ordered; however, the cursor address is not so well ordered. Since we do not have the exact positional information on the cursor, we rely only on CURSOR to properly change the cursor address for the

correct next cursor address depending upon the current position.

The following code is the usual LCD initialization routine for 4-bit interfacing. Here the starting cursor position is position 1 at line 1. The four-bit write routines, `hnibble4`, `instw4`, and `dataw4`, are those we already built and used in Chapter 6. Use them without any change here.

```

;SUBROUTINE LCD4INIT
;Function for 4-bit (only one write must be done)
;In other words, send only the high nibble
LCD4INIT
;IMPORTANT PART
    movlw    0x28
    call    hnibble4
;Function for 4-bit, 2-line display, and 5x8 dot matrix
    movlw    0x28
    call    instw4
;Display On, CURSOR On, No blinking
    movlw    0x0E    ;0F would blink
    call    instw4
;DDRAM address increment by one & cursor shift to right
    movlw    0x06
    call    instw4
;DISPLAY CLEAR
CLEAR
    movlw    0x01
    call    instw4
;
    call    posline1    ;pos1 and line 1
;now CURSOR=1
    return

```

The following subroutine, `LCDisplay`, is the main displaying routine monitoring and handling the cursor positions and addresses.

```

;LCD DISPLAYING SUBROUTINE
LCDisplay
    call    dataw4    ;write a character
    incf    CURSOR    ;every time of display, increase cursor
;CURSOR is automatically incremented by 1 from LCDisplay
;if CURSOR is 20 (0x14), change to posline2
;if CURSOR is 40 (0x28), change to posline3
;if CURSOR is 60 (0x3C), change to posline4
;if CURSOR is 80 (0x50), change to posline1
    clrf    STATUS
    movf    CURSOR,0    ;if the CURSOR is supposed to be
                        ;pos 1 and line 2, the Cursor Address
must
                        ;be changed also
    xorlw   0x15
    btfsc  STATUS, ZERO
    goto   Toline2

    clrf    STATUS
    movf    CURSOR,0

```

```

xorlw      0x29
btfsc     STATUS,ZERO
goto      Toline3

clrf      STATUS
movf      CURSOR,0
xorlw     0x3D
btfsc     STATUS,ZERO
goto      Toline4

clrf      STATUS                ;if the cursor is at the last pos
                                ;at the 4th line, the next cursor
                                ;position is the pos 1 at line 1
                                ;after clearing the LCD

movf      CURSOR,0
xorlw     0x51
btfsc     STATUS,ZERO
call      LCDclearhome          ;delete all an move to (1,1)
return

Toline2
call     posline12
return

Toline3
call     posline13
return

Toline4
call     posline14
return

```

Clearing the LCD and returning to position 1 at line is done by the following subroutine, LCDclearhome. The first two instruction write clears and move the cursor to the "home" position. The next, third, writing is not necessary but used anyway to show the actual cursor address and the variable CURSOR we use throughout our example code.

```

;SUBROUTINE
;DISPLAY CLEAR and Cursor to Home position (line 1, position 1)
LCDclearhome
    movlw    0x01
    call     instw4
;Now let's move the cursor to the home position (position 1 of line #1)
;and set the DDRAM address to 0. This is done by the "return home"
instruction.

    movlw    0x02
    call     instw4
;home position
    movlw    0x80
    call     instw4
    movlw    0x01
    movwf    CURSOR
    return

```

The following four subroutines are for moving the cursor address to the first positions of the four lines, respectively. Note and see the matching hex values for CURSOR and the actual cursor address values that are written by instw4 subroutine.

```

posline11
;Position to pos 1 and line 1
    movlw    0x80
    call    instw4           ;Cursor address for (1,1)
    movlw    0x01
    movwf   CURSOR
    return

posline12                ;pos 1 and line 2
    movlw    0xC0
    call    instw4
    movlw    0x15           ;21
    movwf   CURSOR
    return

posline13                ;pos1 and line3
    movlw    0x94
    call    instw4
    movlw    0x29           ;41
    movwf   CURSOR
    return

posline14                pos 1 and line 4
    movlw    0xD4
    call    instw4
    movlw    0x3D           ;61
    movwf   CURSOR
    return

```

The full example code, without subroutine listings, follows below.

```

;KBD4.asm
;
;NOTE: In this program
; BACK SPACE key is not honored
; CR key is not recognized
;
;
;This program is
;1. To read keys from AT or PS/2 keyboard
;2. to display the key on the 20x4 LCD module by Truly (HD44780 compatible)
;3. Displays from the first dot matrix to the last one
;4. 4. IF the last dot is reached, it is cleared and restart from the first
dot
;
; LCD is with 4-bit interfacing
;

```

```

;CR key would change the line
;
; Pin Connection from LCD to 16F877
; LCD (pin#)      16F877 (pin#)
;DB7 (14) -----RB7(40)
;DB6 (13) -----RB6(39)
;DB5 (12) -----RB5(38)
;DB4 (11) -----RB4(37)
;E (6)  -----RB2(35)
;RW (5)  -----RB3(36)
;RS (4)  -----RB1(24)
;Vo (3)  -----GND
;Vdd (2) -----+5V
;Vss (1) -----GND
;
;KEYBOARD Interfacing
;CLOCK -----RD7 (input)
;DATA -----RD6 (input)
;
;

        list P = 16F877

STATUS      EQU    0x03
PCL         EQU    0x02           ;For Key Table Calling
PCLATH     EQU    0x0A           ;upper part of PC
CARRY      EQU    0x00
ZERO       EQU    0x02
PORTB      EQU    0x06
TRISB      EQU    0x86
RS         EQU    0x01           ;RB1
E          EQU    0x02           ;RB2
RW         EQU    0x03           ;RB3
TRISD      EQU    0x88
PORTD      EQU    0x08
CARRY      EQU    0x00
MSB        EQU    0x07
CLOCK      EQU    0x07           ;from Keyboard (RD7)
KDATA      EQU    0x06           ;from Keyboard (RD6)

;RAM

        CBLOCK    0x20
                CURSOR           ;tracking the current display position
;CURSOR
;1  2  3  4  5  6  7  8  9 10 11 12 13 14 15 16 17 18 19 20  line 1
;21 22 23 24 25 26 27 28 29 30 31 32 33 34 35 36 37 38 39 40  line 2
;41 42 43 44 45 46 47 48 49 50 51 52 53 54 55 56 57 58 59 60  line 3
;61 62 63 64 65 66 67 68 69 70 71 72 73 74 75 76 77 78 79 80  line 4
;
                Daddr           ;Display address (cursor pos)
                Dkey            ;Key character to be displayed
                DATAreg
                Bitcount
                Kstat
                Kount120us      ;Delay count (number of instr cycles for delay)

```

```

        Kount100us
        Kount1ms
        Kount10ms
        Kount1s
        Kount10s
        Kount1m
        Temp          ;temp storage
    ENDC

;program should start from 0005h
;0004h is allocated to interrupt handler

        org          0x0000
        goto        START

Start
        org          0x05

        banksel     TRISD
; 1100 0000

        movlw       B'11000000' ;RB7 for CLOCK and RB6 for DATA as inputs
        movwf      TRISD

        call        delay1s          ;Give Keyboard to send STATUS to the host

        BANKSEL    TRISB
        movlw       0x00
        movwf      TRISB          ;All output

        banksel     PORTB
        clrf       PORTB          ;RW set LOW here

        clrf       CURSOR          ;Current Display Location
        incf       CURSOR          ;Home cursor position (1, 1)
;LCD routine starts
        call        delay10ms
        call        delay10ms

        banksel     PORTB
        clrf       PORTB          ;RW set LOW here
        ;give LCD module to reset automatically
        call        LCD4init

;=====
;KBD Monitoring
BEGIN

        banksel     DATAreg
        clrf       DATAreg

; CHECK IF THE CLCOK is HIGH at least for 10mS

        banksel     PORTD
        btfss      PORTD, CLOCK

```

```

        goto     BEGIN          ;if CLOCK is LOW, start again
        call    Delay10ms      ;10mS delays

;check again for CLCOK
        btfss   PORTD, CLOCK
        goto    BEGIN
;READY FOR CLOCK PULSES

        clr     KSTAT
KEYIN
;X reading
        call    RX11bit        ;
        clr     STATUS
        movf    DATAreg,0     ;Break Code?
        xorlw   0xF0
        btfss   STATUS,ZERO
        goto    CAT
;BREAK is detected. Abort It. Resume It
        goto    BEGIN
;Category detection
CAT     clr     STATUS
        movf    DATAreg,0
        xorlw   0xE0
        btfsc   STATUS,ZERO
        goto    Begin         ;E0 keys (CAT2) are ignored
        clr     STATUS
        movf    DATAreg,0
        xorlw   0x12
        btfsc   STATUS,ZERO
        goto    LRSHIFT
        clr     STATUS
        movf    DATAreg,0
        xorlw   0x59
        btfsc   STATUS,ZERO
        goto    LRSHIFT
        clr     STATUS
        movf    DATAreg,0
        xorlw   0x58         ;CAPS LOCK
        btfsc   STATUS,ZERO
        goto    CAPS
        movf    DATAreg,0
        clr     STATUS          ;CR check
        xorlw   0x5A
        btfsc   STATUS,ZERO
        goto    CRhandle
                                ;L Shift ==>12 | F0 12
                                ;R Shift ==>59 | F0 59
;CAT1 has the format of (X)|(F0)(X)
CAT1   movf    DATAreg,0
;check if the key in is CR
;Then we have to move the next line

        call    NoShiftKeyTable ;(X) display
        call    LCDisplay
;(F0) detection
        call    RX11bit
        clr     STATUS

```

```

    movf      DATAreg,0
    xorlw     0xF0
    btfss    STATUS,ZERO
;Key is not broken. Still pressed,
    goto     CAT1
;Key is broken
;Last (X) reading
    call     RX11bit          ;(X) after F0

    goto     EGIN

;L-SHIFT and R-SHIFT has the form
;L-SHIFT and a character 12 X | F0 X |F0 12
;R-SHIFT and a character 59 X | F0 X |F0 59

LRSHIFT          ;12 or 59 entered

;(F0) detection
    call     RX11bit
    clrf     STATUS
    movf     DATAreg,0
    xorlw     0xF0
    btfsc    STATUS,ZERO
    goto     BEGIN

                                ;X

    clrf     STATUS          ;if (12) do not display
    movf     DATAreg,0
    xorlw     0x12
    btfsc    STATUS, ZERO
    goto     LRSHIFT

    clrf     STATUS          ;if (59) do not display
    movf     DATAreg,0
    xorlw     0x59
    btfsc    STATUS, ZERO
    goto     LRSHIFT

                                ;a Key (X) is entered
    movf     DATAreg,0
    call     ShiftKeyTable
    call     LCDisplay

;(F0) detection
    call     RX11bit
    clrf     STATUS
    movf     DATAreg,0
    xorlw     0xF0
    btfss    STATUS,ZERO
    goto     LRSHIFT
;Last (X) reading
    call     RX11bit
    movf     DATAreg,0
    clrf     STATUS          ;check if (X) or (12) entered after F0
    xorlw     0x12

```

```

        btfsc     STATUS,ZERO
        goto     BEGIN
        goto     LRSHIFT
;
CAPS                                ;caps lock (58) entered

;(F0) detection
        call     RX11bit             ;this must be F0
        call     RX11bit             ;this must be (58) again

CAPnext
        call     RX11bit             ;Check if (58) or other
        clrf     STATUS
        movf     DATAreg,0
        xorlw    0x58
        btfss   STATUS,ZERO
        goto     CAPtwo              ;End of CAP session
        call     RX11bit             ;F0
        call     RX11bit             ;(58)
        goto     BEGIN

;a Key (X) is entered

CAPtwo
        movf     DATAreg,0
        call     CAPKeyTable
        call     LCDisplay

;(F0) detection
        call     RX11bit             ;this
        clrf     STATUS
        movf     DATAreg,0
        xorlw    0xF0
        btfss   STATUS,ZERO
        goto     CAPtwo

;Last (X) reading                    ;F0 is read
        call     RX11bit             ;(X) again and ignore
        goto     CAPnext

;CR handling
CRhandle
        call     RX11bit             ;F0 read
        call     RX11bit             ;CR reading again
;have to move the cursor to the next line at the first position
;
;Routine Here

        goto     BEGIN

;SUBROUTINES and TABLES HERE
;
;HERE
        END
;END of CODE

```

6. A complete Keyboard-LCD Operation with Back Space and Carriage Return

The next version of Keyboard and LCD connection with a 16F877 is to accommodate the two keys we did not: Back Space (BS) and Carriage Return (CR) keys. When BS key is detected we have to move the cursor position to the left by 1 position. This, then, needs the current cursor position address. When CR is detected, again, the cursor position address must be provided so that we move the cursor to the first position of the next line.

These new features require reading information from the LCD controller/interface inside the LCD module, which we have not discussed at all. We only discussed about writing instructions and data to the LCD controller. The reason we need reading information, especially the current cursor address information, is that, as we mentioned before, we do not have direct knowledge on the cursor position address unless we monitor the position every time we write a character on the LCD. Of course there is an indirect way to fulfill this task: monitoring CURSOR value every time we write a character to LCD and interpret the cursor position address from the CURSOR value.

As discussed above, CURSOR is the variable assigned for convenience; it does not indicate the cursor position inside the DD(Display Data) RAM of the LCD module. But reading the cursor position address would be better because we need not this reading all the time as we have to do in the CURSOR tracing approach; we do this cursor position address reading only when CR or BS key is detected. In conclusion, whether we use an arbitrary variable for cursor tracking or not, reading the location address from the LCD interface/controller brings more convenience and gives a good practice of more utilizing a LCD module.

As we studied in Chapter 6, there is a command in HD44780 or equivalent LCD controller/interface of "Reading Busy Flag and DD RAM address" as shown below. On additional command we include is the cursor address setting command. These two are most relevant in the discussion of this version of keyboard - LCD connection.

Instruction	Code										Description	
	RS	R/W	DB7	DB6	DB5	DB4	DB3	DB2	DB1	DB0		
Set DDRAM address	0	0	1	ADD	Sets DDRAM address. DDRAM data is sent and received after this setting.							
Ready busy flag & address	0	1	BF	AC	Reads busy flag (BF) indicating internal operation is being performed and reads address counter contents.							

ADD: DDRAM address (corresponds to cursor address)

AC: Address counter used for DDRAM address

As the table shows, we read 7-bit DDRAM address of the current cursor position. We do not care much on the flag bit, BF, if the LCD is ready to receive data or not, since our interest is the 7 bits returned to 16F877.

By the way, cursor address read from the LCD module for 20x4 display format based on the position and the line is as follows:

Position	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20
First line	00h	01h	02h	03h	04h	05h	06h	07h	08h	09h	0Ah	0Bh	0Ch	0Dh	0Eh	0Fh	10h	11h	12h	13h
Second line	40h	41h	42h	43h	44h	45h	46h	47h	48h	49h	4Ah	4Bh	4Ch	4Dh	4Eh	4Fh	50h	51h	52h	53h
Third line	14h	15h	16h	17h	18h	19h	1Ah	1Bh	1Ch	1Dh	1Eh	1Fh	20h	21h	22h	23h	24h	25h	26h	27h
Fourth Line	54h	55h	56h	57h	58h	59h	5Ah	5Bh	5Ch	5Dh	5Eh	5Fh	60h	61h	62h	63h	64h	65h	66h	67h

To read the cursor address, we have to clear the RS pin and set the RW bit. Since we use 4-bit interfacing, we need two consecutive commanding to the LCD to read the total 8-bit data composed of the BF bit and the 7-bit cursor address. Care should be taken when we read cursor address from LCD module. Usually we set all the pins of PORTB as outputs since the upper 4 data lines RB<7:4> and RW, E, RS are all outputs. However, when we read the data bits RB<7:4> must be changed to input pins.

Another caution we have to use is that in both readings, even though there are only 4 data lines involved, when we read data from PORTB as a byte oriented instruction we read the whole 8 bits anyway. Since the first reading gets upper nibble and the second, the lower nibble, we have to extract the nibbles properly to form a byte hex number. The following code illustrates the DDRAM address reading subroutine, `readad4`, in 4-bit interfacing environment.

As you see at the bottom of the subroutine, once the reading is done, we move back to the usual writing mode by setting the PORTB<7:4> as outputs and clearing the RW line.

```

;subroutine reading the cursor position
;RW Must be High
;RS Must be Low
;the 7th bit is BF flag (so ignore this one, or make MSB 0)
;PORTB <7:4> as inputs
;High then Low nibbles of ADDRESS
;The content of DDADDR read from LCD module (HEX Numbers)
;Line 1: 00 01 02 ..... 13
;Line 2: 40 41 42 ..... 53
;Line 3: 14 15 16 ..... 27
;Line 4: 54 55 56 ..... 67

readad4
    banksel    TRISB                ;set Rb7 - DR4 as inputs
    movlw     0xF0                  ;upper 4 bits as inputs
    movwf     TRISB
    banksel    PORTB
    bsf       PORTB, RW             ;READING MODE
    call      delay1ms
    bcf       PORTB, RS
    call      delay1ms
    bsf       PORTB, E
    call      delay1ms
    bcf       PORTB, E              ;Reading starts here now
                                        ;upper byte first

```

```

;now PORTB<7:4> has the DDRAM ADDRESS
;upper nibble

movlw      0xF0
andwf     PORTB,0
movwf     DDtemp1      ;Ddtemp1 = DDADDR<7:4>|0000
;Reading for the second nibble

bcf       PORTB,RS
call     delay1ms
bsf      PORTB, E
call     delay1ms
bcf      PORTB, E

;reading starts now
;for lower byte
;PORTB<7:4> has the DDADDR<3:0>

movlw      0xF0
andwf     PORTB,0
movwf     DDtemp2      ;Ddtemp2 = DDADDR<3:0>|0000
swapf    DDtemp2      ;Ddtemp2 = 0000|DDADDR<3:0>

;add DDtemp1 and DDtemp2 for DDADDR
;
movf      DDtemp1,0
addwf    DDtemp2,0
movwf    DDADDR      ;DDADDR=DDADDR<7:4>|DDADDR<3:0>

;END of Reading
;Change to Write Mode
banksel  TRISB
movlw    0x00
movwf    TRISB      ;all outputs again
banksel  PORTB
bcf      PORTB,RW   ;back to writing mode
return

```

The next subject is to include in the Category classification section of the code BS and CR key detection part. This is done by simply adding a few lines of the entered key check. The addition for these two keys is at the bottom of this category classification part of the code.

```

;Category detection
CAT   clrf      STATUS
      movf     DATAreg,0
      xorlw    0xE0
      btfsc   STATUS,ZERO
      goto    Begin      ;E0 keys (CAT2) are ignored
;Shift Key Detection

      clrf     STATUS
      movf    DATAreg,0
      xorlw   0x12
      btfsc   STATUS,ZERO
      goto    LRSHIFT
      clrf    STATUS
      movf    DATAreg,0
      xorlw   0x59
      btfsc   STATUS,ZERO
      goto    LRSHIFT

```

```

    clrf      STATUS
    movf     DATAreg,0
    xorlw   0x58           ;CAPS LOCK
    btfsc   STATUS,ZERO
    goto    CAPS

    movf     DATAreg,0
    clrf     STATUS       ;CR check
    xorlw   0x5A
    btfsc   STATUS,ZERO
    goto    CRhandle

    movf     DATAreg,0
    clrf     STATUS
    xorlw   0x66
    btfsc   STATUS,ZERO
    goto    BShandle     ;Back Space Handling

```

So when the BS is entered we handle the case by executing the Bshandle part. CR would jump to CRhandle. Let's discuss about handling when CR key is entered. When CR key is entered, we read the next two frames (i.e, F0k and 0Dh break codes) from the keyboard and ignore them, then we change the new cursor position to the first position at the next line.

Therefore we have to know the current cursor position stored in the DDRAM address. Once the cursor position is read, we have to figure out at which line the cursor is located. The cursor position is read in 7-bit format, however, when we set the cursor position by writing an instruction (See above code table), the DB7 pin must be High, so as soon as we read the cursor position address, we set the 7th bit (MSB) of the address, so that we directly write the address as the new cursor position.

The CRhandle routine listed below shows how to find at which line the current cursor is located by the DDRAM address read and to change the new cursor address to the first position of the next line. In the routine, we notice that we use Borrow flag (which is same as the Carry flag used in add instruction) in sub instruction to find the current line position of the cursor, by employing sublw k instruction. The sublw k instruction is to have the operation of (k – W → W; subtract W from k and store the result to W) and check if k is bigger than W or not: if k is bigger than W there is no Borrow so the Borrow flag in STATUS register is set. The Borrow flag is kind of active low flag which clears when the condition is met. Conclusively, if the Borrow bit is set, the k is bigger than W.

```

;CR handling
CRhandle
    call     RX11bit      ;F0 read
    call     RX11bit      ;CR reading again
;read the current cursor position
    call     readad4
;DDADDR has the content
;NOTE: MSB must be 1 in the cursor command
    bsf     DDADDR, MSB   ;set the 7th bit
;if DDADDR<94, then new cursor position is C0
;if DDADDR<E8, then 80

```

```

;if DDADDR<C0, then D4
;if DDADDR<D4, then 94
    clrf        STATUS
    movf        DDADDR,0
    sublw       0x94                ;k-W -->W
    btfsc       STATUS,Borrow      ;No borrow means that k>W
    goto        CR94                ;is less than 94 i.e.,cursor is at line 1
    clrf        STATUS
    movf        DDADDR,0
    sublw       0xC0
    btfsc       STATUS,Borrow
    goto        CRC0                ;cursor in at line 3

    clrf        STATUS
    movf        DDADDR,0
    sublw       0xD4
    btfsc       STATUS,Borrow
    goto        CRD4                ;cursor is at line 2

    clrf        STATUS
    movf        DDADDR,0
    sublw       0xE8
    btfsc       STATUS,Borrow
    goto        CRE8                ;cursor is at line 4
    goto        BEGIN              ;cursor position out of range

CR94 call        posline12          ;move the cursor to pos 1 line 2
    goto        begin
CRC0 call        posline14          ;move the cursor to pos 1 line 4
    goto        BEGIN
CRD4 call        posline13          ;move the cursor to pos 1 line 3
    goto        BEGIN
CRE8 call        posline11          ;move the cursor to pos 1 line 1
    goto        BEGIN

```

The BShandle routine is not very different from CRhandle part. Once BS is entered, the next two frames must read but ignored. What it does is just to reduce the cursor address by 1 and write a cursor position instruction to LCD so that the cursor is 1 place left shifted.

One glitch in this simple procedure is that, when the current cursor is at the first position of a line, the new cursor should be moved to the last position of the line one above. Except that if the cursor is located at the first position of line 1, there is no change and keep the current position even though BS action. Therefore, here again is where the current position is also an important part of the routine.

The majority of the routine is, therefore, dedicated to find if the current cursor is at the first position in any of 4 lines. If the cursor is not at the first position, we simply decrease the cursor address by 1 and write it back for the new DDRAM address. The following routine is for the BS key handler.

```

;BS Handling
BShandle
    movf        DATAreg,0        ;W holds $66

```

```

        call        RX11bit        ;F0 read
        call        RX11bit        ;BS break code
;read the current cursor position
        call        readad4
;DDADDR has the content
; SO move the current to the left
;NOTE: MSB must be 1 for commanding of the cursor position
        bsf        DDADDR, MSB
;if DDADDR = 94, then new cursor position is D3
;if DDADDR = C0, then new position is 93
;if DDADDR = D4, then new position is A7
;if DDADDR = 80, then new position is 80 (NO CHANGE)
; all other cases, new position is (DDADDR - 1)
        clrf      STATUS
        movf      DDADDR,0
        xorlw    0x94
        btfsc    STATUS, ZERO
        goto     DD94                ;cursor in pos 1 line 3
        clrf      STATUS
        movf      DDADDR,0
        xorlw    0xC0
        btfsc    STATUS,ZERO
        goto     DDC0                ;cursor in pos 1 line 2
        clrf      STATUS
        movf      DDADDR,0
        xorlw    0xD4                ;cursor in pos 1 line 4
        btfsc    STATUS,ZERO
        goto     DDD4
        clrf      STATUS
        movf      DDADDR,0
        xorlw    0x80
        btfsc    STATUS,ZERO
        goto     DD80                ;cursor in pos 1 line 1
;all others
        decf      DDADDR
        decf      CURSOR
        movf      DDADDR,0
        call     instw4
        goto     BEGIN

DD94   movlw     0xD3                ;move cursor to pos 20 line 2
        decf      CURSOR
        call     instw4
        goto     BEGIN

DDC0   movlw     0x93                ;move cursor to pos 20 line 1
        decf      CURSOR
        call     instw4
        goto     BEGIN

DDD4   movlw     0xA7                ;move cursor to pos 20 line 3
        decf      CURSOR
        call     instw4
        goto     BEGIN

DD80   movlw     0x80                ;move cursor to pos 1 line 1
        call     instw4

```

```
goto      BEGIN
```

Except these new routines and subroutines, the code for this updated for CR and BS handling is not different from the previous version. Make a full code and run the program to see if the CR and BS keys are working as we intended them to work.